

# VyZX

## Formal Verification of a Graphical Language

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# The ZX Calculus...

... is a graphical language for reasoning about quantum systems

ZX diagrams are open graphs consisting of green “Z” or red “X” “spiders” and “connections” between them



The Z and X spider

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ZX diagrams are **open graphs** consisting of **green** “Z” or **red** “X” “spiders” and “connections” between them

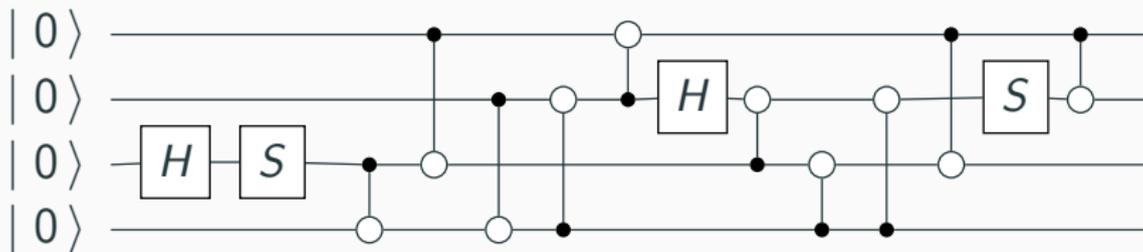


The Z and X spider

Used for **compilation, simulation, error correction** & more

Key benefit: **Diagrammatic rewrites** complete & more comprehensible than circuits or matrices

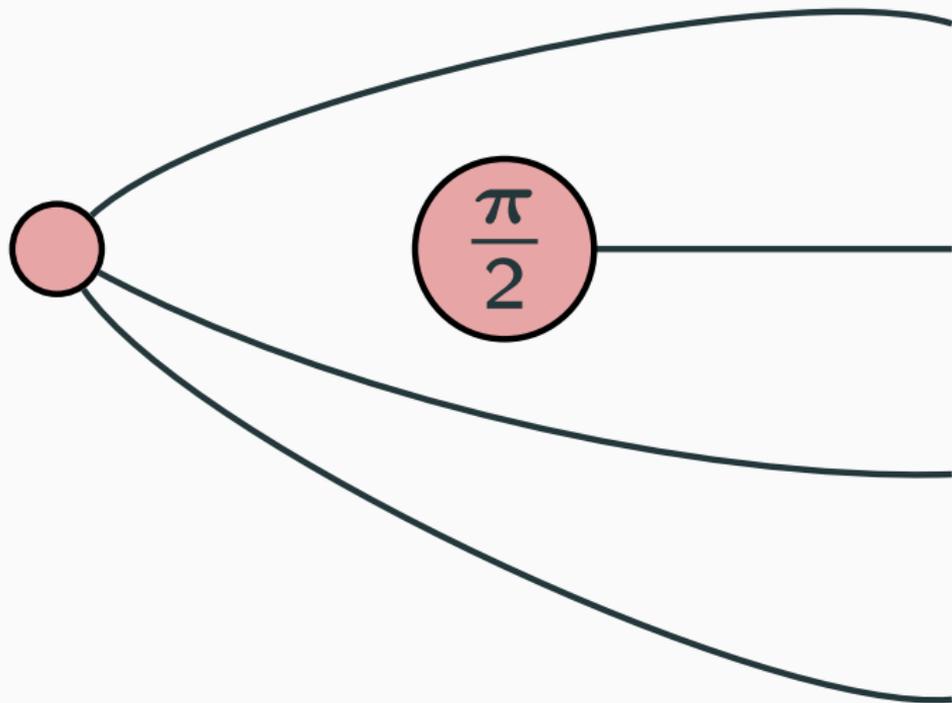
## Example: Entanglement (Van de Wetering)



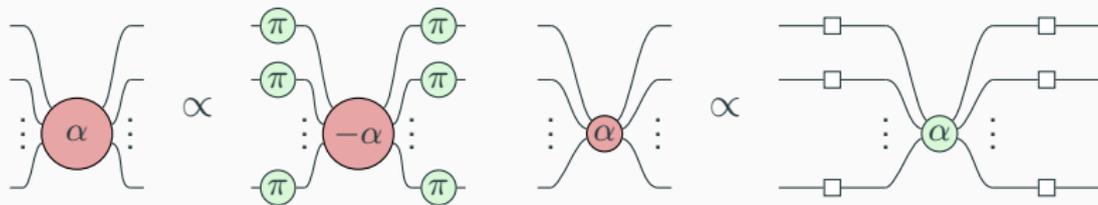
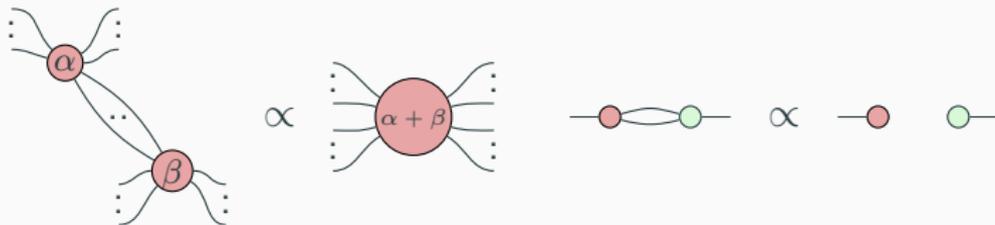
## Example: Entanglement (Van de Wetering)

$$[1 \ 0 \ 0 \ 0 \ i \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1 \ 0 \ 0 \ 0 \ i]^T$$

## Example: Entanglement (Van de Wetering)



# Rewriting Diagrams



Spider Fusion, The Hopf Rule, the Bi-pi Rule (Pi-Copy), Bi-hadamard Rule, the Bialgebra rule, and the identity rule

# Only Connectivity Matters

We can **freely move spiders around**, as long as their connections and in/outputs remain the same



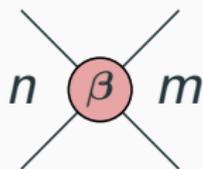
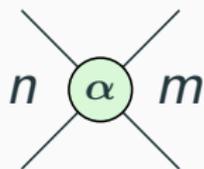
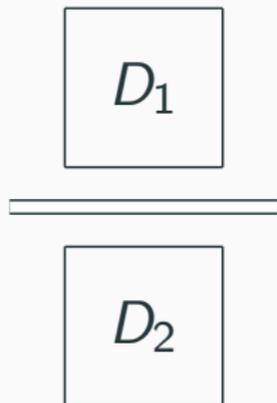
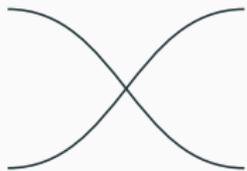
Unverified / Fully axiomatic:

- Quantomatic (<https://quantomatic.github.io/>)
- ZX Calculator ([zx.cduck.me](https://zx.cduck.me))
- Chyp

“Verify” by property testing:

- PyZX (<https://github.com/Quantomatic/pyzx>)

# ZX Diagrams as string diagrams



# Inductive ZX Diagrams

To define our ZX diagrams, we take these string diagram constructions and **add Z and X spiders**.

$$\begin{array}{c} \frac{\text{in out} : \mathbb{N} \quad \alpha : \mathbb{R}}{\text{Z\_Spider in out } \alpha : \text{ZX in out}} \qquad \frac{\text{in out} : \mathbb{N} \quad \alpha : \mathbb{R}}{\text{X\_Spider in out } \alpha : \text{ZX in out}} \\ \\ \text{Cap} : \text{ZX } 0 \ 2 \qquad \text{Cup} : \text{ZX } 2 \ 0 \qquad \text{Swap} : \text{ZX } 2 \ 2 \qquad \text{Empty} : \text{ZX } 0 \ 0 \\ \\ \frac{\text{zx1} : \text{ZX in mid} \quad \text{zx2} : \text{ZX mid out}}{\text{Compose zx1 zx2} : \text{ZX in out}} \qquad \text{Wire} : \text{ZX } 1 \ 1 \\ \\ \frac{\text{zx1} : \text{ZX in1 out1} \quad \text{zx2} : \text{ZX in2 out2}}{\text{Stack zx1 zx2} : \text{ZX (in1 + in2) (out1 + out2)}} \qquad \text{Box} : \text{ZX } 1 \ 1 \end{array}$$

To verify transformations on diagrams, we introduce a system of semantics. Our semantics system will rely on QuantumLib.

$$\begin{aligned} \text{Z\_Spider } n \ m \ \alpha &\mapsto \begin{bmatrix} 1 & \cdots & 0 \\ \vdots & \ddots & 0 \\ 0 & 0 & e^{i\alpha} \end{bmatrix} \\ \text{X\_Spider } n \ m \ \alpha &\mapsto H^{\otimes m} \times \begin{bmatrix} 1 & \cdots & 0 \\ \vdots & \ddots & 0 \\ 0 & 0 & e^{i\alpha} \end{bmatrix} \times H^{\otimes n} \end{aligned}$$

## More Semantics

$$\text{Cap} \mapsto \begin{bmatrix} 1 & 0 & 0 & 1 \end{bmatrix}$$

$$\text{Cup} \mapsto \begin{bmatrix} 1 & 0 & 0 & 1 \end{bmatrix}^T$$

$$\text{Swap} \mapsto \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

$$\text{Empty} \mapsto \begin{bmatrix} 1 \end{bmatrix}$$

$$\text{Wire} \mapsto I_{2 \times 2}$$

$$\text{Box} \mapsto H$$

Compose  $zx1 \ zx2 \mapsto \text{semantics}(zx2) \times \text{semantics}(zx1)$

Stack  $zx1 \ zx2 \mapsto \text{semantics}(zx1) \otimes \text{semantics}(zx2)$

# Proportionality

Equivalence in ZX is up to constant factor

We define **proportionality** and use symbol  $\propto$ :

$$\exists c \neq 0 : \text{semantics}(zx1) = c * \text{semantics}(zx2) \implies zx1 \propto zx2$$

Allows Coq's **rewriting capabilities** in our proofs about diagrams

# Why semantics?

- Smaller TCB
- Interoperability

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Convert circuit components

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- Smaller TCB
- Interoperability
  - We can ingest quantum circuits using sqir
  - Circuit structure is very different
  - Convert circuit components
  - Prove equivalence through ground truth
- $VyZX$  can calculate results

# Three Proof Strategies

## 1. Proof through semantics



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## 2. Inductive proof



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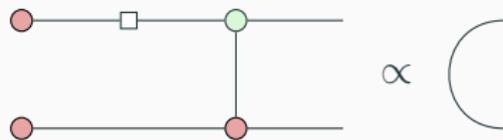
## 1. Proof through semantics



## 2. Inductive proof



## 3. Diagrammatic proof

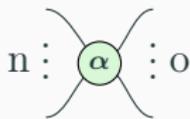


# Three Proof Strategies

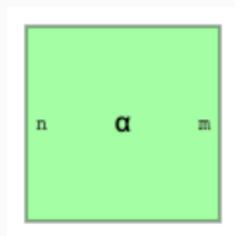
## 2. Inductive proof

$$n \vdots \alpha \quad \vdots m \quad \beta \quad \vdots o \quad \propto \quad n \vdots \alpha + \beta \quad \vdots o$$

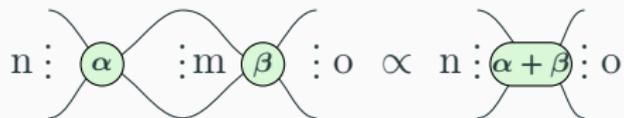
## Other representation



$\cong$

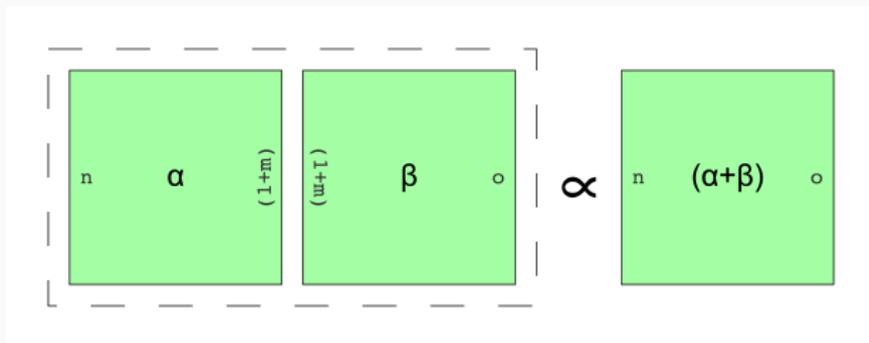


## 2. Inductive proof: Absolute Fusion



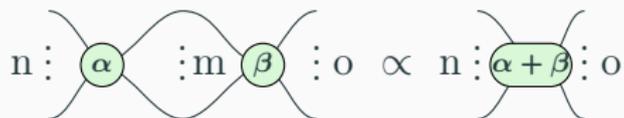
The diagram shows an equality between two configurations of wires and nodes. On the left, there are two nodes,  $\alpha$  and  $\beta$ , each in a green circle. Node  $\alpha$  has  $n$  input wires on the left and  $m$  output wires on the right. Node  $\beta$  has  $m$  input wires on the left and  $o$  output wires on the right. The  $m$  output wires of  $\alpha$  are connected to the  $m$  input wires of  $\beta$ . On the right, there is a single node  $\alpha + \beta$  in a green oval. It has  $n$  input wires on the left and  $o$  output wires on the right. The two configurations are separated by a symbol  $\propto$ .

The lemma:

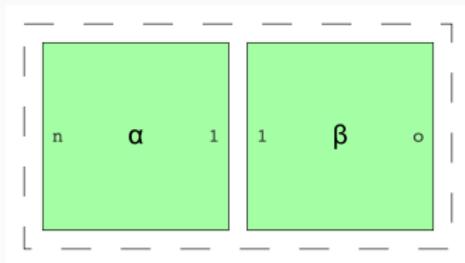


Let's induct over  $m$

## 2. Inductive proof: Absolute Fusion

$$n : \alpha : m : \beta : o \propto n : \alpha + \beta : o$$


Base case:

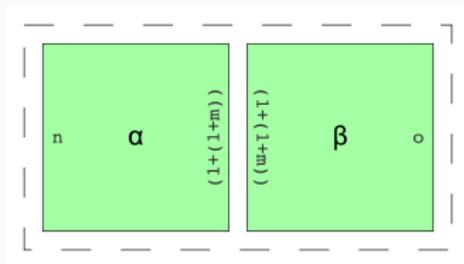


Solve with matrix semantics!

## 2. Inductive proof: Absolute Fusion

$$n \vdots \alpha \quad \vdots m \quad \beta \quad \vdots o \quad \propto \quad n \vdots \alpha + \beta \quad \vdots o$$

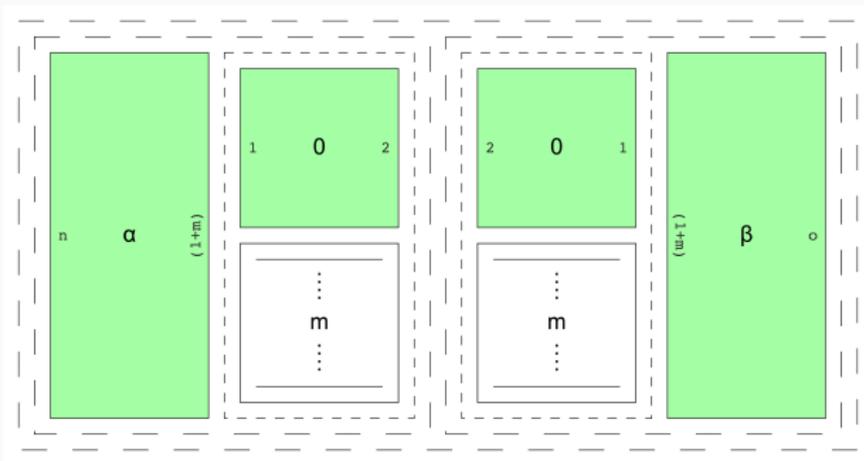
Inductive step:



Let's split out nodes to reduce size

## 2. Inductive proof: Absolute Fusion

$$n \text{ : } \alpha \text{ : } m \text{ : } \beta \text{ : } o \propto n \text{ : } \alpha + \beta \text{ : } o$$

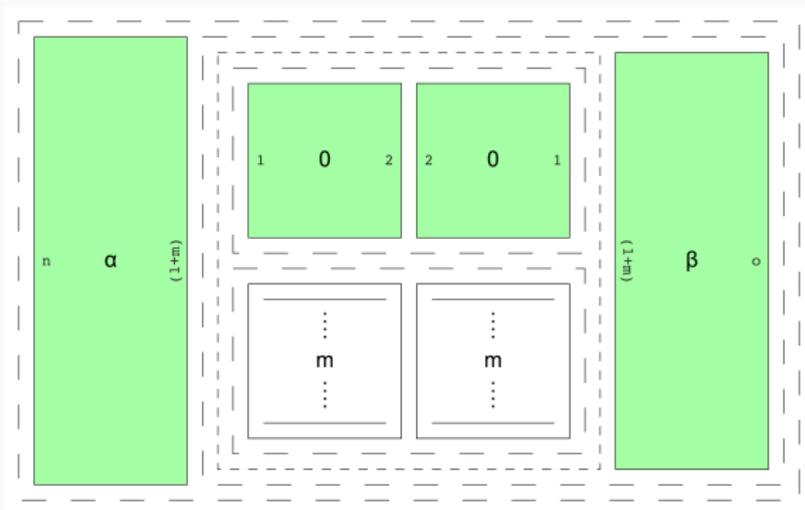


Idea: Fuse the small spiders. Problem: Association

## 2. Inductive proof: Absolute Fusion

$$n \text{ : } \alpha \text{ : } m \text{ : } \beta \text{ : } o \quad \propto \quad n \text{ : } \alpha + \beta \text{ : } o$$

Reassociated diagram

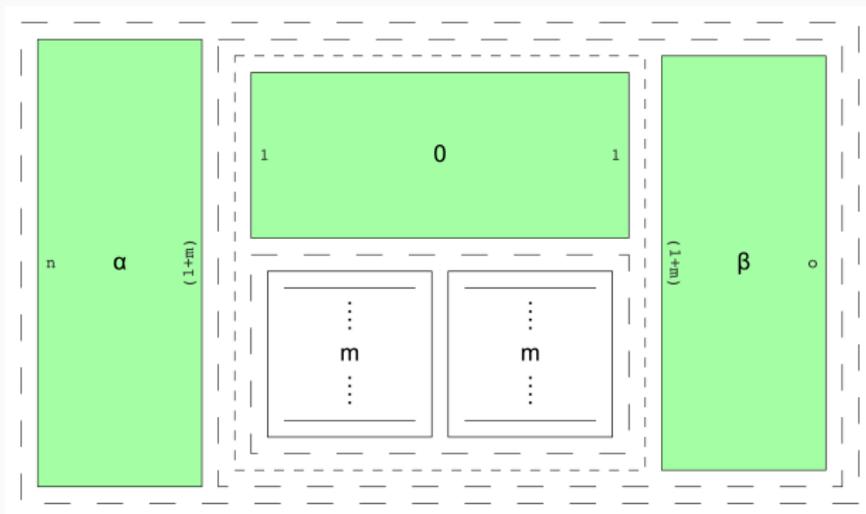


Let's fuse the small spiders

## 2. Inductive proof: Absolute Fusion

$$n \text{ : } \alpha \text{ : } m \text{ : } \beta \text{ : } o \quad \propto \quad n \text{ : } \alpha + \beta \text{ : } o$$

Upper spiders are fused



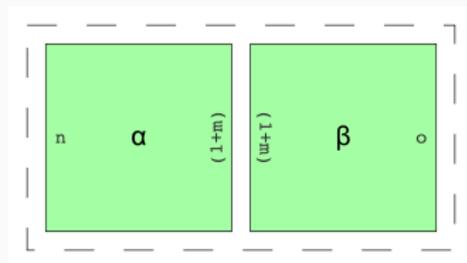
Apply identity rule, remove wires

## 2. Inductive proof: Absolute Fusion

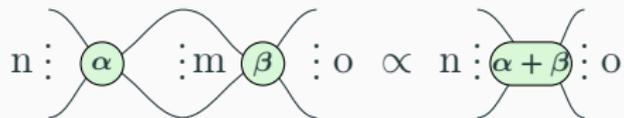
$$n \vdots \alpha \vdots m \vdots \beta \vdots o \propto n \vdots \alpha + \beta \vdots o$$

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Diagram now corresponds to the IH

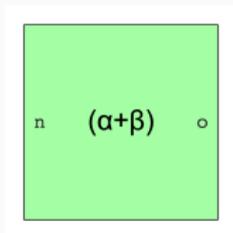


## 2. Inductive proof: Absolute Fusion

$$n \vdots \alpha \vdots m \vdots \beta \vdots o \propto n \vdots \alpha + \beta \vdots o$$


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Qed.



- Any ZX diagram can be expressed
- Multiple ways to encode
- Deal with associativity information
- Dimensionality issues

## How can I verify my graphical language?

- Find underlying categorical structure
- Formally extend structure
- Translate into proof assistant
- Deal with resulting associativity issues

## Future work

- Restore connection information
- Verify ZX-based compiler
- Prove ZX results

# Summary

- Defined ZX diagrams inductively
- Inspired by string diagrams
- Multiple proof strategies
- ZX calculus is interesting!

**Find VyZX on GitHub**

<https://github.com/inQWIRE/VyZX>

**arXiv**

Coming soon...

## References

-  Bob Coecke and Aleks Kissinger, *Picturing quantum processes: A first course in quantum theory and diagrammatic reasoning*, Cambridge University Press, 2017.
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-  Kesha Hietala, Robert Rand, Shih-Han Hung, Xiaodi Wu, and Michael Hicks, *A verified optimizer for quantum circuits*, Proc. ACM Program. Lang. **5** (2021), no. POPL.
-  John van de Wetering, *Zx-calculus for the working quantum computer scientist*, 2020.